Endocentric Constructions and the Cocke Parsing Logic*

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Methods are presented within the parsing logic formulated by Cocke to reduce the large number of intermediate constructions produced and stored during the parsing of even moderately long sentences. A method is given for the elimination of duplicate construction codes stored for endocentric phrases of different lengths.

Automatic sentence-structure determination is greatly simplified if, through the intervention of a parsing logic, the grammatical rules that determine the structure are partially disengaged from the computer routines that apply them. Some earlier parsing programs analyzed sentences with routines that branched according to the grammatical properties or signals encountered at particular points in the sentence, thus having the routines themselves serve as the rules. This not only required separate programs for each language but led to extreme proliferation in the routines, requiring extensive rewriting and debugging with every discovery and incorporation of a new grammatical feature. More recently, programs for sentence-structure determination have employed generalized parsing logics, applicable to different languages and providing primarily for an exhaustive and systematic application of a set of rules.¹⁻⁴ The rules themselves can be changed without changing the routines that apply them, and the routines consequently take fuller advantage of the speed with which digital computers can repeat the same sequence of instructions again and again, changing only the values of some parameters at each cycle.

The case in point is the parsing logic devised by John Cocke in 1960 for applying the rules of a context-free phrase-structure grammar, requiring that each structure recognized by the grammar be analyzed into two and only two immediate constituents (IC).¹

Although all phrase-structure grammars appear to be inadequate in some important respects to the task of handling natural language, they still form the base of the more powerful transformational grammars, which are not yet automated for sentence-structure determination. Moreover, even their severest critic acknowledges that "the PSG [phrase-structure grammar] conception of grammar ... is a quite reasonable theory of natural language which unquestionably formalizes many actual properties of human language" (reference 5, p. 78). Both theoretically and empirically the development and automatic application of phrase-structure grammars are of interest to linguists.

The phrase-structure grammar on which the Cocke parsing logic operates is essentially a table of constructions. Its rules have three entries, one for the code (a descriptor) of the construction, the other two specifying the codes of the ordered pair of immediate constituents out of which it may be formed. The logic iterates in five nested loops, controlled by three simple parameters and two codes supplied by the grammar. They are: (1) the string length, starting with length 2, of the segment being tested for constructional status; (2) the position of the first word in the tested string; (3) the length of the first constituent; (4) the codes of the first constituent; and (5) the codes of the second constituent (Fig.1).

After a dictionary-lookup routine has assigned grammar codes to all the word occurrences in the sentence or total string to be parsed (it need not be a sentence), the parsing logic operates to offer the codes of pairs of adjacent segments to a parsing routine that tests their connectability by looking them up in the stored table of constructions, that is, in the grammar. If the ordered pair is matched by a pair of IC's in the table, the code of the construction formed by the IC's is added to the list of codes to be offered for testing when iterations are performed on longer strings. This interaction between a parsing logic and a routine for testing the connectability of two items is described in somewhat greater detail in Hays.²

In the RAND program for parsing English, the routines produce a labeled binary-branching tree for every complete structural analysis. There will be one tree if the grammar recognizes the string as well formed and syntactically unambiguous and more than one if it is recognized as ambiguous. Even if no complete analysis is made of the whole string, a resume lists all constructions found in the process, including those that failed of inclusion in larger constructions.^{6,7}

Besides simplifying the problem of revising the grammar by separating it from the problem of applica-

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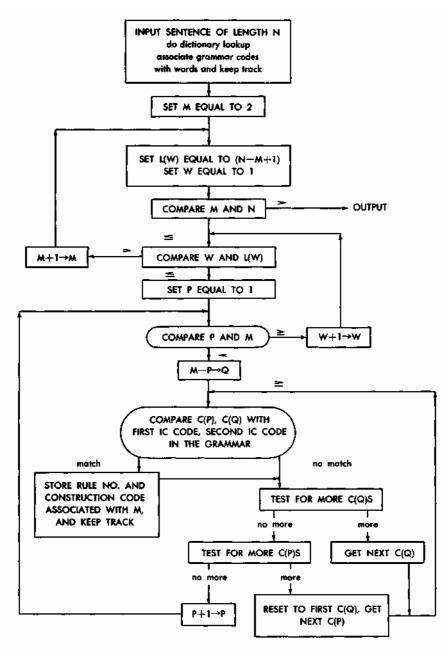


FIG. 1.—Flowchart for the Cocke parsing logic. N: sentence; length; P: string length of first constituent; Q: string length of second constituent; M: P + Q = string length of construction; w: number of first word of construction; L(W): N-M+1 = limit of first word; C(P); code of first constituent; C(Q); code of second constituent.

tion to sentences, the parsing logic, because it leads to an exhaustive application of the rules, permits a rigorous evaluation of the grammar's ability to assign structures to sentences and also reveals many unsuspected yet genuine ambiguities in those sentences.⁸ But because of the difficulties inherent in specifying a sufficiently discriminatory set of rules for sentences of any natural language and because of the very many syntactic ambiguities resolvable only through larger context, this method of parsing produces a long list of intermediate constructions for sentences of even modest length, and this in turn raises a storage problem.

By way of illustration, consider a string of four word occurrences, $x_1 x_2 x_3 x_4$, a dictionary that assigns a single grammar code to each, and a grammar that assigns a unique construction code to every different combination of adjacent segments. Given such a grammar, as in Table 1, the steps in its application to the string by the parsing routines operating with the Cocke parsing logic are represented in Table 2. (The preliminary dictionary lookup assigning the original codes to the occurrences is treated as equivalent to iterating with the parameter for string length set to 1).

TABLE 1 Example Grammar

Rule*	JCI †	1C2‡	cc‡	Rule	IC	ю2	œ
1.	A	в	Е	8.	A]	L
2.	в	С	F	9.	A	ĸ	м
3.	С	D	G	10.	E	G	N
4.	A	F	н	11.	н	D	0
5.	Е	С	I	12.	I	D	Р
6.	в	G	I	13.	А	С	Q
7.	F	D	ĸ	14.		etc.	_

Rule identification tag.

+ Code of first constituent.

Of course, reasonable grammars do not provide for combining every possible pair of adjacent segments into a construction, and in actual practice the growth of the construction list is reduced by failure to find the two codes presented by the parsing logic, when the grammar is consulted. If rule 1 is omitted from the grammar in Table 1, then steps 5, 9, 14, and 16 will disappear from Table 2, and both storage requirements and processing time will be cut down. One method of reducing storage requirements and processing time is to increase the discriminatory power of the grammar through refining the codes so that the first occurrence must belong to class Aa and the second to class Bb whenever adjacent constituents form a construction.

Another way of limiting the growth of the stored constructions is to take advantage of the fact that in actual grammars two or more different pairs of constituents sometimes combine to produce the "same" construction. Assume that A and F (Table 1) combine

 TABLE 2
 Steps Involved in Parsing-Example String*

 (Boxed Section Represents the Parsing-Logic Iterations)

Step	м	w	P	C(P)	c(q)	с(м)	Rule	Steps Combined	Structure Assigned
1.	1	1	1)	Dictionary lookup	x1
2.	1	2	1	В		в	{	Encuonary Rockup	Xa
3.	1	3	1	С		С	٦ ا	assigning codes	X8
4.	1	4	1	D		D	J	to:	X4
4. 5.	2	1	1	A	В	Е	1.	1 + 2	$(x_1 + x_8)$
6.	2	2	1	в	с	F	2.	2 + 3	$(x_2 + x_3)$
6. 7.	2	3	1	С	D	G	3.	3 + 4	$(x_3 + x_4)$
8.	3	1	1		F	н	4.	1+6	$(x_1(x_2+x_3))$
8. 9,	3	1	2	E	С	I	5.	5 + 3	$((x_1 + x_2)x_3)$
10.	3	2	1	в	G	T	6.	2 + 7	$(x_3(x_3 + x_4))$
11.	3	2	2	F	D	ĸ	7.	6 + 4	$((x_{2}+x_{3})x_{4})$
12.	4	1	1	Ā	Ţ	L	8.	1 + 10	$(x_1(x_2(x_3+x_4))))$
13.	4	ī	ī	A	ĸ	M	9.	1 + 11	$(x_1((x_2+x_3)x_4)))$
14.	4	ī	2	E	G	N	10.	$\overline{5} + \overline{7}$	$((x_1 + x_4)(x_3 + x_4)$
15.	$\overline{4}$	ī	3	н	D	0	11.	8+4	$((x_1(x_2+x_3))x_4)$
16,	4	ī	3	I	D	P	12.	9 + 4	$(((x_1 + x_2)x_3)x_4)$

⁹ M: string length of segment; P: length of first-constituent string; C(P): code of first-constituent; C(Q): code of second-constituent string; C(M): code for string, to be stored when C(P) and C(Q) are matched in the grammar.

With such a grammar, the number of constructions to be stored and processed through each cycle increases in proportion to the cube of the number of words in the sentence. If the dictionary and grammar assign more than one code to occurrences and constructions, the number may grow multiplicatively, making the storage problem still more acute. For example, if x_1 were assigned two codes instead of one, additional steps would be required for every string in which x_1 was an element, and iteration on string-length 4 would require twice as many cycles and twice as much storage. to form a construction whose syntactic properties are the same, at least within the discriminatory powers of the grammar, as those of the construction formed by E and c. Then rules 4 and 5 can assign the same code, H, to their constructions. In consequence, at both step 8 and step 9 in the parsing (Table 2), H will be stored as the construction code C(M) for the string $x_1 x_2 x_3$ even though two substructures are recorded for it, that is, $(x_1(x_2 + x_3))$ and $((x_1 + x_2)x_3)$. The string can be marked as having more than one structure, but in subsequent iterations on string-length 4, only one concatenation of the string with x_4 need be made, and step 16 can be omitted. When the parsing has terminated, all substructures of completed analyses are recoverable, including those of marked strings.

Eliminating duplicate codes for the same string from the cycles of the parsing logic results in dramatic savings in time and storage, partly because the elimination of any step has a cumulative effect, as demonstrated previously. In addition, opportunities to eliminate duplicates arise frequently, in English at least, because of the frequent occurrence of endocentric constructions, constructions whose syntactic properties are largely the same as those of one of their elementsthe head. In English, noun phrases are typically endocentric, and when a noun head is flanked by attributives as in a phrase consisting of article, noun, prepositional phrase (A, N, PP), the requirement that constructions have only two IC's promotes the assignment of two structures, (A(N + PP)) and ((A + N)PP), unless the grammar has been carefully formulated to avoid it. Since NP's of this type are common, occurring as subjects, objects of verbs, and objects of prepositions, duplicate codes for them are likely to occur at several points in a sentence.

Consideration of endocentric constructions, however, raises other questions, some theoretical and some practical, suggesting modification of the grammar and the parsing routines in order to represent the language more accurately or in order to save storage, or both. Theoretically, the problem is the overstructuring of noun phrases by the insistence on two IC's and the doubtful propriety of permitting more than one way of structuring them. Practically, the problem is the elimination of duplicate construction codes stored for endocentric phrases when the codes are repeated for different string lengths.

Consider the noun-phrase subject in "All the old men on the corner stared." Its syntactic properties are essentially the same as that of *men*. Fifteen other phrases, all made up from the same elements but varying in length, also have the same properties. They are shown in Table 3.

A reasonably good grammar should provide for the recognition of all sixteen phrases. This is not to say that sixteen separate rules are required, although this would be one way of doing it. Minimally, the grammar must provide two rules for an endocentric NP, one to combine the head noun or the string containing it with a preceding attributive and another to combine it with a following attributive. The codes for all the resulting constructions may be the same, but even so, the longest phrase will receive four different structural assignments or bracketings as its adjacent elements are gathered together in pairs, namely:

(all (the (old (men (on the corner))))), (all (the ((old men) (on the corner)))), (all ((the (old men)) (on the corner))), ((all (the (old men))) (on the corner)).

TABLE 3	TA	BL	Е	3
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	Length				Nou	m P	hras	e	
1.	7	All	the	old	men	on	the	corner	(stared)
2.	6							corner	
3.	6	All	the		men	on	the	corner	
3. 4. 5.		All		old	men	on	the	corner	
5.	5			Old	men	on	the	corner	
6.	6 5 5		The		men	on	the	corner	
7.	5	All			men	on	the	corner	
8.	4				Men	on	the	corner	
9.	4	All	the	old	men				
10.	3		The	old	men				
11.	3 3	All	the	+	men				
12.	3	All		old	men				
13.	2				men				
14.	2		The		men				
15.	2	All			men				
16.	ī				Men				

If it is assumed that the same code, say that of a plural NP, has been assigned at each string length, it is true that only one additional step is needed to concatenate the string with the following verb when the parsing-logic iteration is performed for string-length 9. But meanwhile a number of intermediate codes have been stored during iterations on string lengths 5, 6, 7, and 8 as the position of the first word of the tested string was advanced, so that the list also contains codes for:

men on the corner stared	(length 5),
old men on the corner stared	(length 6),
the old men on the corner stared	(length 7),
Il the old men on the corner stared	(length 8).

а

Again, the codes may be the same, but duplicate codes will not be eliminated from processing if they are associated with different strings, and strings of different length are treated as wholly different by the parsing logic, regardless of overlap. If this kind of duplication is to be reduced or avoided, a different procedure is required from that available for the case of simple duplication over the same string.

But first a theoretical question must be decided. Is the noun phrase, as exemplified above, perhaps really ambiguous four-ways, and do the four different bracketings correlate systematically with four distinct interpretations or assignments of semantic structure?⁸ And if so, is it desirable to eliminate them? It is possible to argue that some of the different bracketings do correspond to different meanings or emphases orin earlier transformational terms-to different orderings in the embeddings of "the men were old" and "the men were on the corner" into "all the men stared." Admittedly the native speaker can indicate contrasts in meaning by his intonation, emphasizing in one reading that all the men stared and in another that it was all the old men who stared; and the writer can resort to italics. But it seems reasonable to

assume that there is a normal intonation for the unmarked and unemphatic phrase and that its interpretation is structurally unambiguous. In the absence of italics and other indications, it seems unreasonable to produce four different bracketings at every encounter with an NP of the kind exemplified.

One way to reduce the duplication is to write the grammar codes so that, with the addition of each possible element, the noun head is assigned a different construction code whose distribution as a constituent in larger constructions is carefully limited. For the sake of simplicity, assume that the elements of NP'S have codes that reflect, in part, their ordering within the phrase and that the NP codes themselves reflect the properties of the noun head in first position and are subsequently differentiated by codes in later positions that correspond to those of the attributes. Let the codes for the elements be 1 (all), 2 (the), 3 (old), 4 (men), 5 (on the corner). Rules may be written to restrict the combinations, as shown in Table 4. With these rules, the grammar provides for only one structural assignment to the string:

(all (the (old (men + on the corner)))).

This method has the advantage of acknowledging the general endocentricity of the NP while allowing for its limitations, so that where the subtler differences among NP'S are not relevant, they can be ignored by ignoring certain positions of the codes, and where they are relevant, the full codes are available. The method should lend itself quite well to code-matching routines for connectability. However, if carried out fully and consistently, it greatly increases the length and complexity of both the codes and the rules, and this may also be a source of problems in storage and processing time.²

Another method is to make use of a classification of the rules themselves. Since the lowest loop of the parsing logic (see Fig. 1) iterates on the codes of the second constituents, the rules against which the paired strings are tested are stored as ordered by first IC codes and subordered by second IC codes. If the iterations of the logic were ordered differently, the rules would also be ordered differently for efficiency in testing. In other words, the code of one constituent in the test locates a block of rules within which matches for all the codes of the other constituent are to be sought; but the hierarchy of ordering by one constituent or the other is a matter of choice so long as it is the same for the parsing logic and for storing the table of rules that constitute the grammar. In writing and revising the rules, however, it proves humanly easier if they are grouped according to construction types. Accordingly, all endocentric NP's in the RAND grammar are given rule identification tags with an N in first position. Within this grouping, it is natural to subclass the rules according to whether they attach attributives on the right

or on the left of the noun head. If properly formalized, this practice can lead to a reduction in the multiple analyses of NP's with fewer rules and simpler codes than those of the previous method.

TABLE 4

Rule	ıcl	ю2	cc	
1.	1+	4	→ 4 1	(All men)
2.	2 +	4	\rightarrow 42	(The men)
3.	3 +	4	→ 43	(Old men)
4. 5.	4 +	5	$\rightarrow 45$	(Men on the corner)
5.	1 +	42	→ 412	(All the men)
6.	1+	43	→ 413	(All old men)
7.	2 +	43	→ 423	(The old men)
8.	ĩ÷	423		(All the old men)
9.	1 ∔		→ 415	(All men on the corner); but not *41 + 5 \rightarrow 415
10.	2 +	45	$\rightarrow 425$	(The men on the corner); but not $*42 + 5 \rightarrow 425$
11.	3 +	45	→ 4 35	(Old men on the corner); but not *43 + 5 \rightarrow 435
12.	2 +	435	→ 4235	(The old men on the corner); but not *423 + $5 \rightarrow 4235$
13.	1 +	4235	→ 41235	(All the old men on the corner); but not °4123 + 5 \rightarrow 41235

As applied to the example, the thirteen rules and five-place codes of Table 4 can be reduced to two rules with one-place codes and an additional feature in the rule identification tag. The rules can be written as:

Although the construction codes are less finely differentiated, the analysis of the example will still be unique, and the number of abortive intermediate constructions will be reduced. To achieve this effect, the connectability-test routine must include a comparison of the rule tag associated with each C(P) and the rule tags of the grammar. If a rule of type *N is associated with the C(P), that is, if an *N rule assigned the construction code to the string P which is now being tested as a possible first constituent, then no rule of type \$N can be used in the current test. For all such rules, there will be an automatic "no match" without checking the second constituent codes (see Fig. 1). As a consequence of this restriction, in the final analysis, the noun head will have been combined with all attributives on the right before acquiring any on the left.

To be sure, the resume of intermediate constructions will contain codes for "old men," "the old men," and "all the old men," produced in the course of iterations on string lengths 2, 3, and 4, but only one structure is finally assigned to the whole phrase, and the intermediate duplications of codes for strings of increasing length will be fewer because of the hiatus at stringlength 5. For the larger constructions in which the NP participates, the reduction in the number of stored intermediate constructions will be even greater.

Provisions may be made in the rules for attaching still other attributives to the head of the NP without great increase in complexity of rules or multiplication of structural analyses. Rule \$N2, for example, could include provision for attaching a relative clause as well as a prepositional phrase, and while a phrase like "the men on the corner who were sad" might receive two analyses unless the codes were sufficiently differentiated to prevent the clause from being attached to *corner* as well as to *men*, at least the further differentiation of the codes need not also be multiplied in order to prevent the multiple analyses arising from endocentricity.

Similarly, for verb phrases where the rules must allow for an indefinite number of adverbial modifiers, a single analysis can be obtained by marking the strings and the rules and forcing a combination in a single direction. In short, although the Cocke parsing logic tends to promote multiple analysis of unambiguous or trivially ambiguous endocentric phrases, at the same time increasing the problem of storing intermediate constructions, the number of analyses can be greatly reduced and the storage problem greatly alleviated if the rules of the grammar recognize endocentricity wherever possible and if they are classified so that rules for endocentric constructions are marked as left (*) or right (\$), and their order of application is specified.

A final theoretical-practical consideration can at least be touched on, although it is not possible to develop it adequately here. The foregoing description provided for combining a head with its attributives (or dependents) on the right before combining it with those on the left, but either course is possible. Which is preferable depends on the type of construction and on the language generally. If Yngve's hypothesis9 that languages are essentially asymmetrical, tending toward right-branching constructions to avoid overloading the memory, is correct, then the requirement to combine first on the right is preferable. This is a purely grammatical consideration, however, and does not affect the procedure sketched above, in principle. For example, consider an endocentric construction of string-length 6 with the head at position 3, so that its extension is predominantly to the right, thus: 1 2 (3) 4 5 6. If all combinations were allowed by the rules, there would be thirty-four analyses. If combination is restricted to either direction, left or right, the number of analyses is reduced to eleven. However, if the Cocke parsing logic is used to analyze a left-branching language, making it preferable to specify prior combination on the left, then the order of nesting of the fourth and fifth loops of the parsing logic should be reversed (Fig. 1) and the rules of the grammar should be stored in order of their second constituent codes, subordered on those of the first constituents.

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