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**ORGANIZATION AND PROGRAMMING  
OF THE  
MULTISTORE PARSER**

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The Multistore system was developed in order to recognize and explain structural patterns in natural-language sentences (specifically English) and eventually yield an output in which the relations between the various items of the sentence are hierarchically displayed.

The recognition of these structural patterns is made by means of a system of rules which operate on a sequence of words, i.e. a sentence, whose individual characteristics are pre-established. By individual characteristics are meant the possibilities a word has to correlate (i.e. to form a syntactic combination) with another item; these possibilities are represented by 'correlators', that is, by syntactic elements which link two items in a correlation.

Each word is characterized by a set of pre-established data:

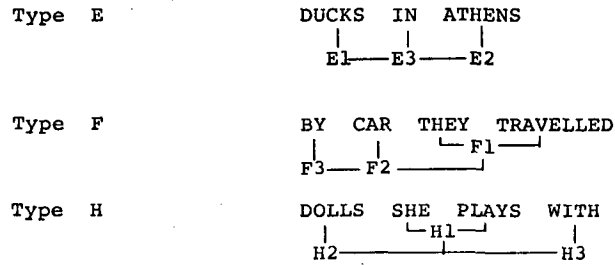
- a) the S-code, which distinguishes between the various senses of a homograph. For instance, a word like "READ" will have four different S's to distinguish between:

READ = supine	e.g. I CAN READ
READ = past tense	YESTERDAY I READ
READ = past part.	I HAVE READ
READ = noun	A LONG READ

These distinctions are essential, since whenever a homograph occurs, one and only one of its meanings can be taken into consideration to make the final pattern, unless, of course, the sentence is ambiguous and more than



EXPLICIT CORRELATOR



For each type there are different correlational functions which determine the position a word has in a correlation. When two adjacent words have complementary functions of the same Ic - for instance, word A has 5050 N<sub>1</sub> and word B has 5050 N<sub>2</sub> - a 'product' is made and recorded in the form:

Word A 5050 N Word B

This product is considered as one piece and can become first or second correlatum in a wider correlation and is therefore treated as though it were a single word, i.e., it is assigned strings of Ic's which indicate its correlational possibilities both with adjacent words and with adjacent products already made. Single words, however, being vocabulary items, have their strings of Ic's assigned a priori; products, since they arise during the procedure, have to be assigned their Ic-strings dynamically. The assignation of specific Ic's to a product depends on:

- a) the correlator responsible for the particular correlation;
- b) the characteristics (Ic's) of the word (or product) which

makes up the first or the second correlatum.

The operational cycle that assigns Ic's to a product we call 'reclassification'.

The amount of data involved in an analysis of this kind is really enormous. Let us consider a sentence consisting of ten words, each of which has two different senses (S's). On an average 50 correlational indices are assigned to each sense of a word. Now, just to check the correlational compatibility of two adjacent words about 10,000 matching operations would be necessary; the matching procedure for all the words of the sentence would involve about 90,000 operations. On an average five products would result from the first 10,000 matching operations; each of them would be assigned about 50 correlational indices that represent the product's correlational possibilities to correlate with another adjacent piece - either a word or a product. The procedure to match these five products with another piece would involve about 637,000 operations. If to this figures we add the number of operations necessary starting from level 3 (see p. 7) with all the products made in the immediately preceding levels (200,000), the total number of operations involved would come to 927,000.

The reclassification routine also involves a great number of operations of this kind: about half of the correlational indices a product is assigned depends on the corre-

C 3 C 3

C 55 C 5  
P 0007 01 0 /

P 0011

S 0014

lator responsible for that correlation; the other half depends on the strings of indices which the two correlata of the product have. According to the presence or absence of specific indices in the strings of the first or second correlatum, pre-established sets of indices are assigned to the product; or sets of indices are assigned to the product only if they are present in the strings of its two correlata. The reclassification of each product would require about 2,000 operations, which means 100,000 for the average of 50 products in a sentence of 10 words, bringing the total of operations to over a million only for the matching procedure. This would imply - for this part of the program alone - processing times of the order of some seconds of machine time if the most modern computer is available, or of about an hour - at best - if the work is done with an older model.

The amount of work and money involved in a procedure of this kind made us try to find a quicker and more economical way of handling correlational indices: as a result of our efforts the Multistore system was developed. (Bibl. 1)

The basic idea of the Multistore consists in pre-establishing in a given area of the machine's central core as many separate positions as there are correlators in the system. The arrangement of these positions representing the correlators orders them according to type. This assures that at any point of the procedure each Ic is not used sev-





Every column is dedicated to one Ic and subdivided into two subcolumns, if the Ic is of type N, M, or V (implicit); if the Ic is of type E, F, or H (explicit), the column is divided into three subcolumns.

The lines L1, L2, L3 etc. divide the area into levels. The levels are determined by the succession of words in input. Thus each level bears the number of the word it represents. Every input word causes for each Ic in its Ic string the insertion of a marker into the Multistore column corresponding to that Ic; and the level of that marker in the Ic column corresponds to the input number and the position of that word in the sentence. Thus all the markers inserted for one word represent the correlational possibilities of that word.

If, on the line of level 1, an Ic of the string representing correlator type N of the first word has caused the insertion of a marker into the column corresponding to that Ic and if an Ic of string N of the second word has caused the insertion, into the line of level two, of a marker representing CF2 of the same Ic, this implicitly means that with the same correlator a correlation is made containing the first and the second word; this product No x, consisting of word 1 (S 'a') and word 2 (S 'b') is the product of correlator No y and type N. This product belongs to level 2 and when it has been assigned a string of Ic's by the appropriate rules of reclassification, it will be inserted

into the Multistore on the second level; this means that it can enter into combinations only with those words that belong to the immediately preceding level, or with products which contain the words of the immediately preceding level. Such a correlation, whenever it is made, would still belong to the level of product x. In our specific case product x could correlate only with an item of level zero, which does not exist, because product x is on level two and already contains word No. 1. Hence we can formulate a restrictive rule to the effect that a product can be a potential second correlatum in an N correlation only if its lower level is larger than 1. The Multistore system lends itself to the introduction of many such restriction rules.

When on a given level all products that have sprung from the insertion of markers corresponding to the word of that level have been reclassified, and the products originating from that reclassification have, in turn, been reclassified and have inserted their markers, and there are no more products to be reclassified, then the procedure inserts the next word and thus begins the next level. This means that once a subsequent word of the sentence has been inserted, all preceding words and products become 'inactive' pieces, having exhausted every possible attempt of correlation with 'active' pieces; from then on they represent merely latent correlational possibilities with subsequent items. This state of inactivity in the case of the last word of the sen-

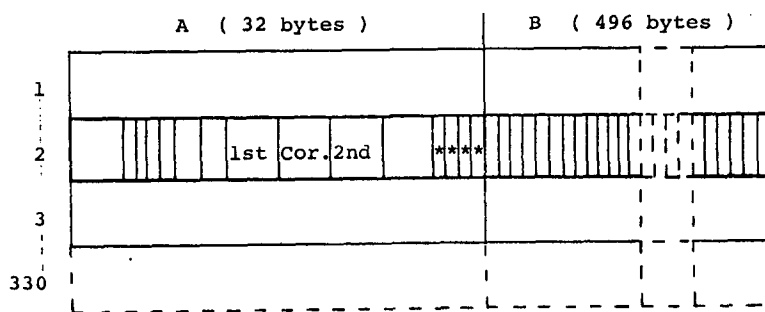
tence determines the end of the analysis. At this point the product (or products) that contains all words of the sentence is called 'complete' and represents the hierarchical structure of the sentence.

The first tentative program MP1 (Bibl. 1) was written for use on a GE 425 computer and its main purpose was to show the applicability of the Multistore system to correlational grammar and to check the method of programming based on 'significant addresses'.

The present program, MP2 (Bibl. 3) , is a revised and enlarged version of MP1 written for use on an IBM 360/67 computer. On the basis of our previous experience it can be considered an actual working tool.

Many solutions, as well as many restraints, depend on the fact that under many respects it is a machine-oriented program. The program is structured on a large area of the central core, divided into lines and columns, whose size is 528 x 330 bytes. Each line (330 all together) consists of 528 bytes and is divided into two sections: A and B. Section A contains all the data necessary to define a line; section B consists of 496 bytes, that is, of as many bytes as there are correlators operative in the system. Each line specifies as permanent data a reclassification rule - whose definition is given in section A of the same line - and the set of indices assigned by that rule (bit 6). The relevance of the rule to a given product

is specified in the columns of section B.



\* reclassification rule

Fig. 2

Each byte of section B is divided into 8 bits as illustrated below.

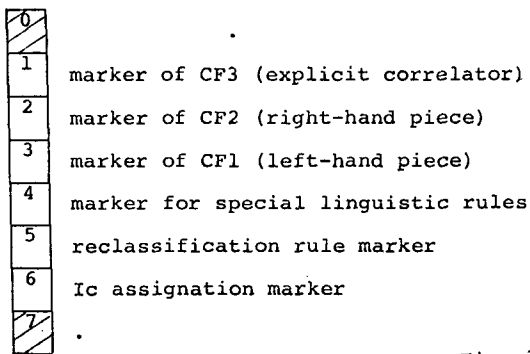


Fig. 3

Bits 1 to 4 are therefore used in the matching procedure, whereas bit 5 and 6 are pre-established data to be used in the reclassification routine.

#### Procedure

Each S-value of a word occupies one line of the Multi-store area and its specifications are recorded in section A of the same line. For each  $I_c$  contained in the string of that S-value of the word, a marker is inserted, according to the correlational function, in bit 1,2 or 3 of the corresponding byte of the line, that is, in the byte which bears that  $I_c$  as label.

According to its function, a marker can be a left-hand piece 'LH', and as such it is simply recorded, or a right-hand piece 'RH', in which case, immediately after it has been recorded, the column is searched for a complementary and contiguous LH piece. If this is found, an indication of product is recorded in the first free line of the Multi-store; this address consists of three data: a) the address of the line where the LH piece was found, which is recorded in the area 'first correlatum'; of the line of the product; b) the address of the line where the RH piece was recorded, which is recorded in the area 'second correlatum' and c) the relative address of the column which characterizes both both LH and RH pieces, which is recorded in the area 'correlator'.

After the product's specifications are recorded, and if there are no other LH pieces with which the RH piece in hand can combine, the routine for the insertion of  $I_c$ 's is resumed. If the insertion of the next  $I_c$  of that S-value



of the word causes a new product to be made, the procedure is repeated and the product is recorded on the next free line of the Multistore area. Only when all the Ic's of the piece which has caused the production have been inserted, the reclassification routine takes place, starting from the first product newly recorded.

The information contained in the area 'correlator' of the line containing the product's record gives the address of the Multistore column dedicated to the correlator responsible for that product. The column is then searched, from the top down, for a bit 5 set ON (see Fig. 3 on p.10). If it is found, this implicitly means that on the line to which the bit belongs, there will be found the record of a reclassification rule relevant to the product to be reclassified. Section A of the same line contains the instructions concerning the assignation of the Ic's whose markers are contained in bit 6 (see Fig.3). A bit 6 set ON implicitly indicates either the column in which to check (if the rule requires it) the record of the first or second correlatum (the addresses of which are recorded in section A of the line in which the rule is recorded) for the presence or the absence (as specified by the rule) of bits 1 to 3 set ON (which represent Ic's); or it may simply indicate the place in which to insert a marker, i.e. a reclassification Ic. The routine for the insertion of reclassification markers is exactly the same as the routine for the insertion of markers for words.

W/P	S	lst.	Corr.	2nd.	Rules	1	2	3	4	5	⑥	7	8	9	0	1	1	1	1	4	
						1	2	3	4	5	6	7	8	9	0	1	2	3	6		
0000	W1	Sa				1															
L1	528	W2	Sa			2															
						3	x	x	x	x											
	1056	W2	Sp			4															
						5															
L2	1584	Px		0000	6	1056	6														
							7														
	2112						8														
							9														
	2640						0														
							1														

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\* Conditioned rule. 1 Unconditioned rule. # Check on 1st correlatum.  
 Ic 4 CF1. @ Assign the string CF1 contained in the rule to the product.  
 [ Assign the string CF2 contained in the rule to the product.

RECLASSIFICATION DIAGRAM



The analysis of the sentence is complete when the last marker of the last word-sense has been inserted and there are no further products to be reclassified or re-cycled. At this point the output routine starts. Three different kinds of output are produced:

- a) a list of all the products made in the course of the analysis of the sentence;
- b) a list of all Ic's assigned to each product during the reclassification routine ;
- c) a graphic representation of the hierarchical structure of all 'complete' products (that is, containing all words of the sentence). This structure is equivalent to a tree structure with words at the terminals and correlators at the nodes. (see Appendix)

This is a general outline of the procedure of combination, production, reclassification and output. In addition to that there are several routines which meet special requirements. A special rule, for instance prevents specific RH pieces from becoming eligible LH pieces once a certain correlation - which contains them as RH pieces - has been made. A word like "LITTLE", for instance, in its function as a quantifier, once it has been correlated with the definite article and made the product "THE//LITTLE" cannot become LH piece in the correlation:

LITTLE // HE KNOWS

The indication 'discard' on print-out type 'a' - i.e. on the list of all the products made during the analysis - will show that "LITTLE" is no more available as LH piece for any other correlation.

Another restraint concerns some 'complete' products which, though grammatically correct, cannot be accepted as interpretation of the sentence. For instance, in a sentence like:

THEY // WERE READY

the structure which takes "WERE" as subjunctive is not acceptable, since it would require something else - an "IF" or "I WISH" etc. - to precede. In cases like this the indication 'non-sentence' appears in print-out type 'a'.

A set of special routines serves the purpose of recognizing idiomatic expressions. When one of them is recognized, inserted in the Multistore and reclassified -like any other product - the indication 'idiom' is printed on print-out type 'a'.

The whole program, including the Multistore area and buffers, is contained in the central core of the machine and occupies about 200 K. The system accepts as input, sen-

tences of up to 16 words - a limit fixed in accordance with the average length of sentences in scientific texts (Bibl.5) and ample enough to allow any type of syntactic structure. Processing-times for 10-word sentences are about 1-1.5 seconds. Our present vocabulary is limited to 150 words for reasons of punched card maintenance. However, it could be enlarged without affecting the program.

The Multistore parser was developed for the automatic analysis of English sentences on the basis of correlational grammar, but it is in no way limited to this kind of grammar. Actually, by changing the input parameters, the symbols of the rules and the matching operations it could be used to handle the data of any kind of predictive grammar; neither are its dimensions critical; the Multistore area could be reduced or enlarged simply by altering the parameters in accordance with the storage capacity available in the machine.

B i b l i o g r a p h y

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- 3 The Multistore System MP-2 (E.v.Glasersfeld and P.P. Pisani), Scientific Progress Report, Grant AFOSR 1319-67 Georgia Institute for Research, Athens, Georgia, 1968.
- 4 The Multistore Parser for Hierarchical Syntactic Structures (E.v.Glasersfeld and P.P.Pisani) Grant AFOSR 1319-67, Georgia Institute for Research, Athens, Georgia 1969 (paper submitted to Communications of ACM)
- 5 Computational Analysis of Present-Day American English (Henry Kucera and W.Nelson Francis), Brown University Press, Providence, Rhode Island, 1967

Input Sentence: C1 I C2 3FAC 03 11 04 FIS 05 FRIGHT 06 80PK  
 07 ON 08 1001AS

(Print-out type a)

P 0001	W 0001 01 01	2220 N	W 0002 01 05	C1 03
P 0002	W 0001 01 01	2250 N	W 0002 04 05	C1 03
P 0003	W 0002 04 05	4010 N	W 0003 01 03	C2 04
P 0004	W 0002 02 05	4010 N	W 0003 01 03	C2 04
P 0005	W 0002 01 05	4010 N	W 0003 01 03	C2 04
P 0006	W 0001 01 01	2250 N	P 0003	C1 04
P 0007	W 0001 01 01	2250 N	P 0005	C1 04
P 0008	W 0004 02 05	5210 N	W 0005 01 15	C4 06
P 0009	P 0009	2220 N	W 0006 01 05	C4 07
P 0010	P 0009	4010 N	W 0003 01 03	C2 07
P 0011	W 0005 01 15	5210 N	W 0006 05 05	C5 07 DISCARD
P 0012	W 0004 01 05	7016 N	P 0011	C4 07
S 0013	P 0012	0550 E	W 0007 01 07	C4 08

APPENDIX I  
 Complete Parsing

S CC15	W CCC6 01 VS	0550 E	W 0007 01 07	C4 09
P CC16	W CCC6 01 VS	4470 N	W 0007 02 04	C4 09
P CC17	F 0008	2220 N	P 0016	C4 09
P CC18	P CC17	4010 N	W 0003 01 47	C3 09
P CC19	W CCC6 01 VS	0550 E	W 0009 02 74	C6 09
P CC20	F CC10	0550 E	W 0008 02 74	C3 09
P CC21	F CC12	0550 E	W 0009 02 74	C4 09
P CC22	F 0008	2220 N	F 0019	C4 09
P CC23	W CCC3 01 F3	5571 V	P 0021	C3 09
P CC24	F 0022	4010 N	W 0003 01 42	C3 09
P CC25	W CCC2 04 FF	4010 N	F 0023	C2 09
P 0026	W CCC2 03 FF	4010 N	P 0022	C2 09
P CC27	W CCC2 01 FS	4010 N	F 0023	C2 09
P 0028	W 0001 01 41	2250 N	P 0025	C1 09 COMPLETE

(Print-out type c)

178540

P C28	2250N
P C25	4010N
P C23	0571V
P C21	0550E
P C12	7016A
P C11	5210A
REF ID: A63, HIS CP PRI 1A RCC 88 CN TCA 7	

P 0025      K C001 01 81      2220 N      P 0027      . C1 00      COMPLETE

178542

P 025	+	-----	2220N-+
P 027	.	+	4010N-+
P 022	.	+	0571V-+
P 021	.	+	0550S-+
P 012	.	+	7016N-+
P 011	.	+	4521CA-+
I	.	REA FS	IV 13, FIS CF ERI 1A ECC 18 CH 1P1 1E.

AI